# Alternatives to block size as aggregate resource limit

Mark Friedenbach <mark@friedenbach.org> Scaling Bitcoin Montréal, QC 2015

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## Background: why have a block size?

- Place a finite upper bound on resources required to validate a Bitcoin block
  - A hard upper bound on size of buffers during block transmission and validation (engineering considerations, primarily)
  - 2) Rate-limit resource consumption during validation (achieve decentralization requirements)
  - 3) Other limits e.g. MAX\_BLOCK\_SIGOPS derived from block size

## Some problems emerge...

- Block size correlates with resource consumption in the typical case
  - But design criteria must be met even for worst-case, adversarial situations.
- Specially constructed blocks can be made that require significantly more resources to validate than a typical 1MB block
  - Observed in practice!
- Actual limit must be constrained by worst-case scenario
  - How much worse is the worst case?
  - Pretty bad, actually...

### F2Pool spam cleanup

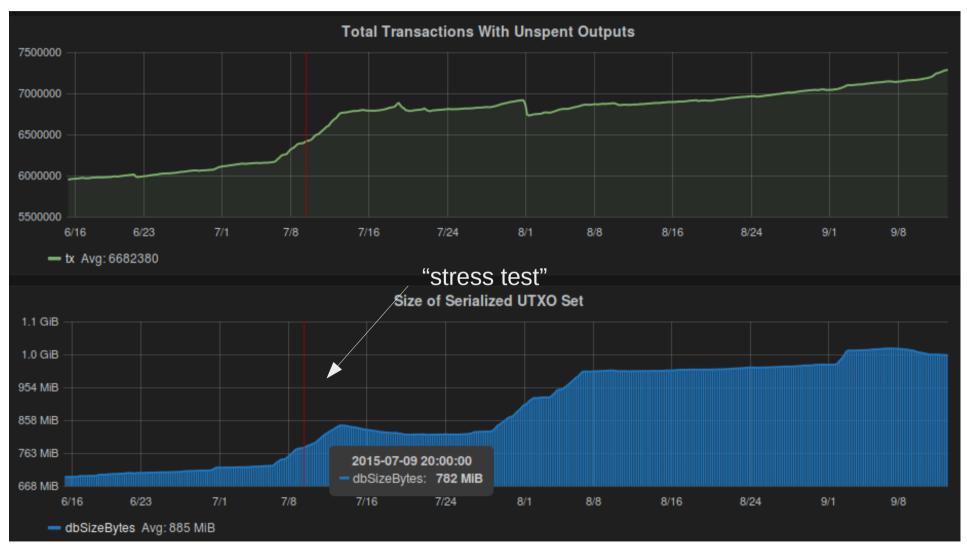
txid:bb41a757f405890fb0f5856228e23b715702d714d59bf2b1feb70d8b2b4e3e08

- Block with only 1 non-coinbase transaction
  - Sweeps 5569 dust UTXOs. Size: 999,657 bytes.
- Transaction re-serialized for each signature check
  - Total 1.25GB of data serialized & hashed.
- ~30s to verify on actual nodes at the time.
- Scales as O(n²)
  - 3.2 MB: 10 min
  - 8.0 MB: 2 hr 8 min

## MAX\_BLOCK\_SIGOPS FAIL CVE 2013-2292

- MAX\_BLOCK\_SIGOPS limits the aggregate number of signature checks in the outputs of a block...
  - ...but it is the inputs, not outputs that are run.
- Vulnerable to attack
  - Over time create outputs with 200 CHECKSIG's each.
  - Spend all in one giant transaction.
  - MAX\_BLOCK\_SIGOPS does not apply.
- "A transaction that takes at least 3 minutes to verify" (Sergio Damian Lerner, 30 Jan 2013) https://bitcointalk.org/?topic=140078

## **UTXO** set growth



http://statoshi.info/dashboard/db/unspent-transaction-output-set

#### How bad is it?

#### Worst case is pretty bad

- Between 10x 100x slowdown from typical
- Attacks are cheap (fees not linked to real costs)
- O(n<sup>2</sup>) scaling gets worse with larger block size
- Attacks observed in the wild!

We need a new measure of resource consumption that tracks validator costs more accurately than block size alone

#### Factors which affect full validation

#### Block size

worst-case latency

#### UTXO growth

- created minus spent

#### Script...

- opcodes executed
- space required
- bytes copied

## Elliptic curve operations

In inputs, not outputs!

## Bytes hashed

Adjusted by algorithm?

#### Bytes copied

- OP\_DUP...

## A linear function of many variables

### Infinite possible functions to consider

- Future work?
- But...

#### A linear combination of factors

- Simplest commitment structure for fraud proofs (Merkle sum tree)
- Straightforward, easy to implement solvers
- Drop-in replacement in existing infrastructure

#### Selection of coefficients

#### Some factors are directly comparable

 Convert opcode execution counts, signature validations, and bytes hashed to single-threaded CPU running time.

#### Type error in some comparisons

- How many bytes of RAM equals 100 ms CPU utilization?
- Use available server hardware to establish conversion ratios.

#### Factors grow differently over time

- Some factors expected to increase with Moore's law (parallel CPU speed).
- Others expected to level out in the near future (global latency)

## **Summary & future work**

- Block size meant to rate-limit validater resource consumption
  - Large resource usage causes propagation delays; delays cause centralization pressures.
- Atypical blocks observed in the wild have widely varying resource usage
  - Block size does poor job of predicting resource utilization & propagation delay in an adversarial environment.
- Linear function of multiple factors ideal replacement for block size metric
  - Simple, drop-in replacement for block size metric
  - Requires future work on finalizing set of factors and coefficients

## Thank you!

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See you in Hong Kong, Dec 2015!